#### Introduction to Computer Networks

**COSC 4377** 

Lecture 22

Spring 2012

April 11, 2012

#### Announcements

- HW10 due this week
- HW11 is out
- Student presentations

#### **HW10**

- Latency measurements
- Plotting latency

Get your user id soon

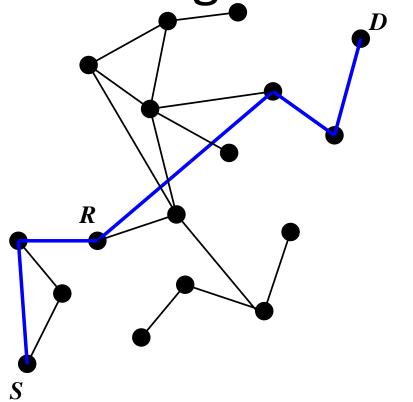
### Today's Topics

- Multi-hop Wireless Networks
- Security
- RPL

### Many Challenges

- Routing
  - Link estimation
- Multihop throughput dropoff

## The Routing Problem



- Find a route from S to D
- Topology can be very dynamic

#### Routing

- Routing in ad-hoc networks has had a lot of research
  - General problem: any-to-any routing
  - Simplified versions: any-to-one (base station),
    one-to-any (dissemination)
- DV too brittle: inconsistencies can cause loops
- DSDV
  - Destination Sequenced Distance Vector

#### **DSDV**

- Charles Perkins (1994)
- Avoid loops by using sequence numbers
  - Each destination increments own sequence number
    - Only use EVEN numbers
  - A node selects a new parent if
    - Newer sequence number or
    - Same sequence number and better route
  - If disconnected, a node increments destination sequence number to next ODD number!
  - No loops (only transient loops)
  - Slow: on some changes, need to wait for root

#### Many Others

- DSR, AODV: on-demand
- Geographic routing: use nodes' physical location and do greedy routing
- Virtual coordinates: derive coordinates from topology, use greedy routing
- Tree-based routing with on-demand shortcuts

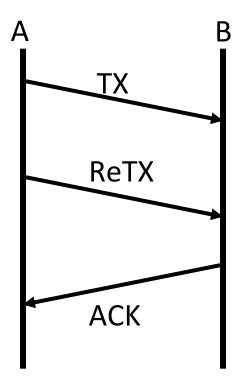
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#### Routing Metrics

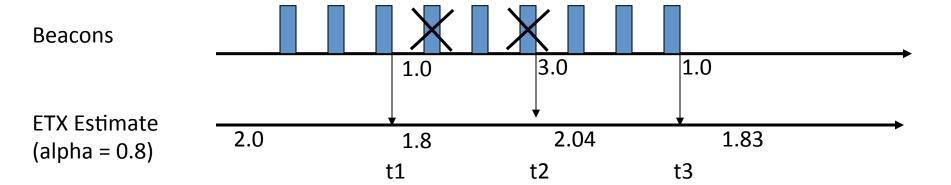
- How to choose between routes?
- Hopcount is a poor metric!
  - Paths with few hops may use long, marginal links
  - Must find a balance
- All links do *local retransmissions*

## **Link Quality Estimation**

ETX(L) = 
$$\frac{1}{PRR(f) * PRR(b)}$$



### **ETX Estimation Example**



#### Routing Metrics

- Idea: use expected transmissions over a link as its cost!
  - -ETX = 1/(PRR) (Packet Reception Rate)
  - Variation: ETT, takes data rate into account

#### Multihop Throughput

1 2 3 6

- Only every third node can transmit!
  - Assuming a node can talk to its immediate neighbors
  - (1) Nodes can't send and receive at the same time
  - (2) Third hop transmission prevents second hop from receiving
  - (3) Worse if you are doing link-local ACKs
- In TCP, problem is worse: data and ACK

# Sometimes you can't (or shouldn't) hide that you are on wireless!

#### TCP over wireless

- How to handle
  - Link losses
  - Hop-by-hop retransmissions
  - Congestion vs lossy links



From: Internal Revenue Service [mailto:admin@irs.gov]

Sent: Wednesday, March 01, 2006 12:45 PM

To: john.doe@jdoe.com

Subject: IRS Notification - Please Read This .



After the last annual calculations of your fiscal activity we have determined that you are eligible to receive a tax refund of \$63.80. Please submit the tax refund request and allow us 6-9 days in order to process it.

A refund can be delayed for a variety of reasons. For example submitting invalid records or applying after the deadline.

To access the form for your tax refund, please click here

Regards, Internal Revenue Service

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## Basic Requirements for Secure Communication

- Availability: Will the network deliver data?
  - Infrastructure compromise, DDoS
- Authentication: Who is this actor?
  - Spoofing, phishing
- Integrity: Do messages arrive in original form?
- Confidentiality: Can adversary read the data?
  - Sniffing, man-in-the-middle
- Provenance: Who is responsible for this data?
  - Forging responses, denying responsibility
  - Not who sent the data, but who created it

#### Other Desirable Security Properties

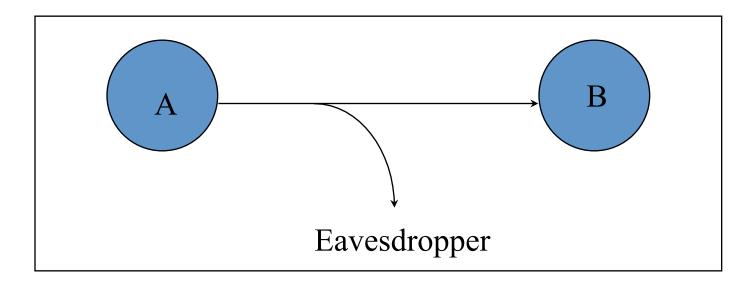
- Authorization: is actor allowed to do this action?
  - Access controls
- Accountability/Attribution: who did this activity?
- Audit/Forensics: what occurred in the past?
  - A broader notion of accountability/attribution
- Appropriate use: is action consistent with policy?
  - E.g., no spam; no games during business hours; etc.
- Freedom from traffic analysis: can someone tell when I am sending and to whom?
- Anonymity: can someone tell I sent this packet?

#### Internet's Design: Insecure

- Designed for simplicity in a naïve era
- "On by default" design
- Readily available zombie machines
- Attacks look like normal traffic
- Internet's federated operation obstructs cooperation for diagnosis/mitigation

## Eavesdropping - Message Interception (Attack on Confidentiality)

- Unauthorized access to information
- Packet sniffers and wiretappers
- Illicit copying of files and programs

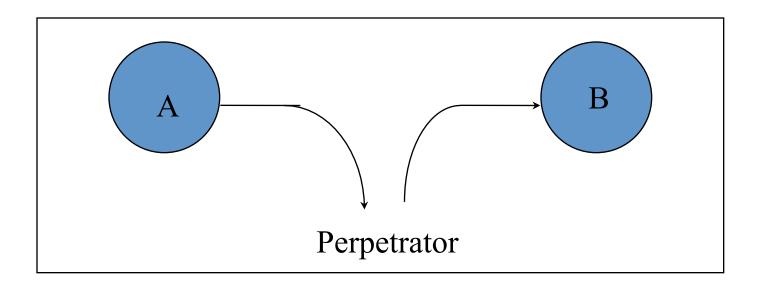


#### Eavesdropping Attack: Example

- tcpdump with promiscuous network interface
  - On a switched network, what can you see?
- What might the following traffic types reveal about communications?
  - DNS lookups (and replies)
  - IP packets without payloads (headers only)
  - Payloads
- How about HW9?

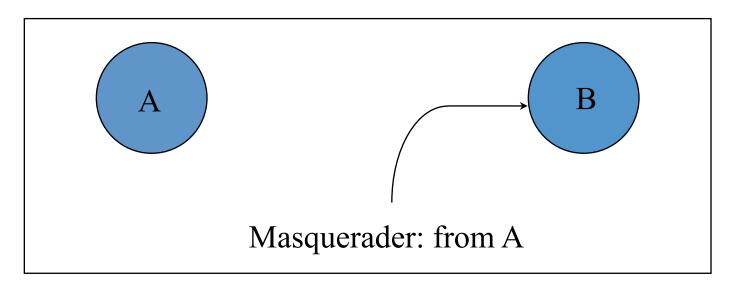
#### Integrity Attack - Tampering

- Stop the flow of the message
- Delay and optionally modify the message
- Release the message again



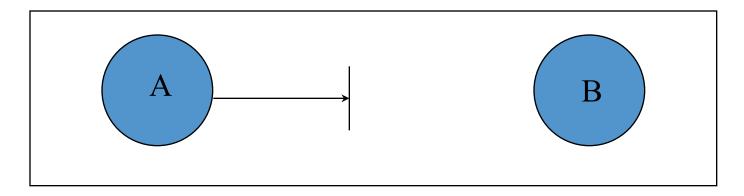
#### **Authenticity Attack - Fabrication**

- Unauthorized assumption of other's identity
- Generate and distribute objects under this identity



#### Attack on Availability

- Destroy hardware (cutting fiber) or software
- Modify software in a subtle way
- Corrupt packets in transit



- Blatant denial of service (DoS):
  - Crashing the server
  - Overwhelm the server (use up its resource)