CHAPTER III
SCHEDULING

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Chapter overview

- The problem

- Non-preemptive policies:
  - FCFS, SJF

- Preemptive policies:
  - Round robin, multilevel queues with feedback, guaranteed scheduling
  - Examples: UNIX, Linux, Windows NT and after
The scheduler

- Part of the OS that decides *how to allocate the processor cores and the main memory* to processes
- Will focus here on the **CPU scheduler**
  - Decides which ready process should get a processor core
  - Also called short-term scheduler
Objectives

- A good scheduler should
  - Minimize *user response times* of all interactive processes
    - *Major objective today*
  - Maximize *system throughput*
  - Be *fair*
  - Avoid *starvation*
What is starvation?

- Starvation happens whenever some ready processes never get core time
  - Typical of schedulers using priorities
    - Lowest-priority processes keep getting set aside
- Remedy is to *increase* the priorities of processes that have waited *too long*
Fairness

- Ensuring fairness is more difficult than avoiding starvation
  
  *If I give freshly-baked cookies to half of my nephews and stale bread to the others, I am not fair but I still ensure that nobody starves*
Non-preemptive Schedulers

- A *non-preemptive* CPU scheduler will never remove a core from a running process
- Will wait until the process releases the core because
  - It issues a system call
  - It terminates
- Now *obsolete*
Examples (I)

- **First-Come First-Served (FCFS):**
  - Simplest and easiest to implement
  - Uses a FIFO queue
  - Seems a good idea but
    - Processes requiring a few ms of core time have to wait behind processes that make much bigger demands
  - *Inacceptable*
How SJF works

- Five students wait for their instructor at the beginning of her office hours
  - Ann needs 20 minutes of her time
  - Bob needs 30 minutes
  - Carol needs 10 minutes
  - Dean needs 5 minutes
  - Emily needs 5 minutes
Examples (II)

- **Shortest Job First (SJF):**
  - Gives a core to the process requesting the least amount of core time
    - Will reduce average wait
    - Must know ahead of time how much core time each process needs
      - *Not possible*
  - Still lets processes monopolize a core
# FCFS schedule

<table>
<thead>
<tr>
<th>Student</th>
<th>Time</th>
<th>Wait</th>
</tr>
</thead>
<tbody>
<tr>
<td>Ann</td>
<td>20</td>
<td>0</td>
</tr>
<tr>
<td>Bob</td>
<td>30</td>
<td>20</td>
</tr>
<tr>
<td>Carol</td>
<td>10</td>
<td>50</td>
</tr>
<tr>
<td>Dean</td>
<td>5</td>
<td>60</td>
</tr>
<tr>
<td>Emily</td>
<td>5</td>
<td>65</td>
</tr>
</tbody>
</table>
The outcome

- Average wait time:
  \[(0 + 20 + 50 + 60 + 65)/5 = 39\text{ minutes}\]
# SJF schedule

<table>
<thead>
<tr>
<th>Student</th>
<th>Time</th>
<th>Wait</th>
</tr>
</thead>
<tbody>
<tr>
<td>Dean</td>
<td>5</td>
<td>0</td>
</tr>
<tr>
<td>Emily</td>
<td>5</td>
<td>5</td>
</tr>
<tr>
<td>Carol</td>
<td>10</td>
<td>10</td>
</tr>
<tr>
<td>Ann</td>
<td>20</td>
<td>20</td>
</tr>
<tr>
<td>Bob</td>
<td>30</td>
<td>40</td>
</tr>
</tbody>
</table>
The outcome

- Average wait time:
  - $(0 + 5 + 10 + 20 + 40)/5 = 15$ minutes

- Less than half the wait time of the FCFS schedule
  - The data were rigged
Preemptive Schedulers

- A *preemptive* scheduler can return a running process to the ready queue whenever another process requires that core in a more urgent fashion:
  - Has been for too long in the ready queue
  - Has higher priority

- *Sole acceptable solution*
  - Prevents processes from “hogging” a core
Types of preemptive schedulers

- *Preemptive schedulers w/o priorities:*
  - All processes have the same priority
  - Ready queue is FIFO

- *Preemptive schedulers with priorities:*
  - Use multiple queues
  - Differ in the way they adjust process priorities
Round robin (I)

- Assumes all processes have \textit{same priority}
  - Guaranteed to be starvation-free
- Similar to FCFS but processes only get the a core for \textit{up to} $T_{CPU}$ time units
  - \textit{Time slice} or \textit{time quantum}
- Processes that exceed their time slice return to the end of the ready queue
Round robin (II)

Process exceeds time slice

Ready queue

System request completion

Core

System call
How RR works

- Assume
  - Single core
  - Time slice is 100ms (reasonable choice)
  - Ready queue contains processes A, B and C

- A gets core at $t = 0$ms
- A releases the core at $t = 24$ms to do an I/O
- B gets core at $t = 24$ms
- A returns to ready queue at $t = 32$ms
- B forced to release the core at $t = 124$ms
Finding the right time slice (I)

- A small time slice means a good response time
  - No process will ever have to wait more than \((n_{\text{ready\_queue}} + 1)T_{\text{CPU}}\) time units
    where \(n_{\text{ready\_queue}}\) is the number of processes already in the ready queue

- A large time slice means a better throughput
  - Fewer context switches
Finding the right time slice (II)

Ideal CPU schedule

True CPU schedule
The problem

- Want to adjust the time slice to guarantee a maximum waiting time in the ready queue
  \[ T_{CPU} = \frac{T_{max}}{(n_{ready\_queue} + 1)} \]
  - Works well as long as system is lightly loaded
  - Produces very small time slices when system is loaded
    - Too much context switch overhead!
An observation

- The throughput of a system using a RR scheduler actually decreases when its workload exceeds some threshold
  - *Rare* among physical systems
  - *Frequent* among systems experiencing congestion

- Freeway throughput actually decreases when its load exceeds some threshold
The solution (I)

- Add *priorities*

- Distinguish among
  - *Interactive processes*
  - *I/O-bound processes*
    - Require small amounts of core time
  - *CPU-bound processes*
    - Require large amounts of core time *(number crunching)*
The solution (II)

- Assign
  - *High priorities* to interactive processes
  - *Medium priorities* to I/O-bound processes
  - *Low priorities* to CPU-bound processes
The solution (III)

- Assign
  - *Smallest time slices* to interactive processes
  - *Medium time slices* to I/O-bound processes
  - *Biggest time slices* to CPU-bound processes

- Allow higher priority processes to steal cores from lower priority processes
The outcome

- Interactive processes will get good response times
- CPU-bound processes will get the CPU
  - Less frequently than with RR
  - For longer periods of time
  - Less context switch overhead
Two problems

- How to assign priority to processes?
  - Process behaviors may change during their execution
- How to avoid starvation?
Multi-Level with Feedback Queues

- Use *dynamic priorities*

- **Reward**
  - Processes that issue system calls
  - Processes that interact with user
  - Processes that have been a long time in the ready queue

- **Penalize**
  - Processes that exceed their time slice
Implementation (I)

- High priority queue
- Medium priority queue
- Low priority queue

CPU

System call
The priority game

- Different systems have different conventions for priorities
  - **0 is highest**
    - Most UNIX systems, Linux
  - **0 is lowest**
    - UNIX System V Release 4 (V.4)
    - Windows NT and after
Implementation (II)

- Time slice increase when priority decreases, say
  - T for high priority processes
  - 2T for medium priority processes
  - 4T for low priority processes
System V.4 scheduler

- Three process classes:
  - Real-time
  - Time-sharing
  - System (for kernel processes)

- Each process class has its own priority levels
  - Real-time processes have highest priority
  - Time-sharing lowest
Real-time processes

- Have *fixed priorities*
  - As in Windows scheduler
- System administrator can define
  - A different *quantum size* \(\text{rt}_\text{quantum}\) for each priority level
Timesharing processes (I)

- Have *variable priorities*
- System administrator can specify the parameters of each priority level
  - Maximum flexibility
  - Maximum risk of making a bad choice

Leaving too many tuning options to the system administrator increases the chances that some options will be poorly selected!
Timesharing processes (II)

- Parameters include
  - Quantum size (ts_quantum)
  - New priority for processes that use their whole CPU quantum (ts_tqexp)
  - New priority for processes returning from blocking state (ts_slpret)
Timesharing processes (III)

- Maximum amount of time a process can remain in the ready queue without having its priority recomputed \((ts_{maxwait})\)

- New priority for processes that have been in the ready queue for \(ts_{maxwait}\) \((ts_{lwait})\)
Example

<table>
<thead>
<tr>
<th>ts_quantum</th>
<th>ts_tqexp</th>
<th>ts_slpret</th>
<th>ts_maxwait</th>
<th>ts_lwait</th>
<th>LEVEL</th>
</tr>
</thead>
<tbody>
<tr>
<td>1000</td>
<td>0</td>
<td>1</td>
<td>50000</td>
<td>1</td>
<td># 0</td>
</tr>
<tr>
<td>500</td>
<td>0</td>
<td>2</td>
<td>20000</td>
<td>2</td>
<td># 1</td>
</tr>
<tr>
<td>200</td>
<td>1</td>
<td>3</td>
<td>10000</td>
<td>3</td>
<td># 2</td>
</tr>
<tr>
<td>100</td>
<td>2</td>
<td>3</td>
<td>10000</td>
<td>3</td>
<td># 3</td>
</tr>
</tbody>
</table>

- System has four priority levels
  - 0 is lowest
  - 3 is highest
- Anything after a pound sign is a comment
New priorities can be

- Rewarding a “good” behavior: `ts_slpret` and `ts_lwait`
- Penalizing CPU “hogs”: `ts_tqexp`
How?

- We *increase* the priority of processes that
  - Have completed a system call
    - They might become less CPU-bound
  - Have waited a long time in the ready queue
    - To prevent starvation

- We *decrease* the priority of processes that
  - Have exhausted their time quantum
    - They might be more CPU-bound
Second example (I)

<table>
<thead>
<tr>
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</tr>
<tr>
<td>500</td>
<td>X</td>
<td>2</td>
<td>20000</td>
<td>2</td>
<td># 1</td>
</tr>
<tr>
<td>200</td>
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</tr>
<tr>
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</tr>
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<td>3</td>
<td>4</td>
<td>10000</td>
<td>Z</td>
<td># 4</td>
</tr>
</tbody>
</table>

- Table now defines five priority levels

- What are the **correct values** for X, Y and Z?
Second example (II)

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<td>10000</td>
<td>Z</td>
<td># 4</td>
</tr>
</tbody>
</table>

- X is the new priority for processes at level 1 that exceed their time quantum
  - Must be lower than current priority
- X=0
Second example (III)

<table>
<thead>
<tr>
<th>#ts_quantum</th>
<th>ts_tqexp</th>
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</tr>
<tr>
<td>100</td>
<td>3</td>
<td>4</td>
<td>10000</td>
<td>Z</td>
<td>4</td>
</tr>
</tbody>
</table>

- Y is the new priority for processes at level 3 that exceed their time quantum
  - Must be higher than current priority
- Y = 4
## Second example (IV)

<table>
<thead>
<tr>
<th>ts_quantum</th>
<th>ts_tqexp</th>
<th>ts_slpret</th>
<th>ts_maxwait</th>
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<td>3</td>
<td>10000</td>
<td>3</td>
<td># 2</td>
</tr>
<tr>
<td>100</td>
<td>2</td>
<td>4</td>
<td>10000</td>
<td>4</td>
<td># 3</td>
</tr>
<tr>
<td>100</td>
<td>3</td>
<td>4</td>
<td>10000</td>
<td>Z</td>
<td># 4</td>
</tr>
</tbody>
</table>

- **Z** is the new priority for processes at level 4 that have waited too long in the ready queue
  - Should be higher than current priority
  - Level 4 already is the highest priority
- **Z** = 4
Second example (V)

<table>
<thead>
<tr>
<th>#ts_quantum</th>
<th>ts_tqexp</th>
<th>ts_slpret</th>
<th>ts_maxwait</th>
<th>ts_lwait</th>
<th>LEVEL</th>
</tr>
</thead>
<tbody>
<tr>
<td>1000</td>
<td>0</td>
<td>1</td>
<td>50000</td>
<td>1</td>
<td># 0</td>
</tr>
<tr>
<td>500</td>
<td>0</td>
<td>2</td>
<td>20000</td>
<td>2</td>
<td># 1</td>
</tr>
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<td>3</td>
<td>10000</td>
<td>3</td>
<td># 2</td>
</tr>
<tr>
<td>100</td>
<td>2</td>
<td>4</td>
<td>10000</td>
<td>4</td>
<td># 3</td>
</tr>
<tr>
<td>100</td>
<td>3</td>
<td>4</td>
<td>10000</td>
<td>4</td>
<td># 4</td>
</tr>
</tbody>
</table>

- Recall that
  - `ts_slpret` and `ts_lwait` reward “good” behaviors
  - `ts_tqexp` penalizes a “bad” one
An exercise

- Fill the missing values

<table>
<thead>
<tr>
<th>#ts_quantum</th>
<th>ts_tqexp</th>
<th>ts_slpret</th>
<th>ts_maxwait</th>
<th>ts_lwait</th>
<th>LEVEL</th>
</tr>
</thead>
<tbody>
<tr>
<td>1000</td>
<td>X</td>
<td>1</td>
<td>50000</td>
<td>1</td>
<td># 0</td>
</tr>
<tr>
<td>500</td>
<td>Y</td>
<td>2</td>
<td>20000</td>
<td>2</td>
<td># 1</td>
</tr>
<tr>
<td>200</td>
<td>1</td>
<td>3</td>
<td>10000</td>
<td>3</td>
<td># 2</td>
</tr>
<tr>
<td>100</td>
<td>2</td>
<td>Z</td>
<td>10000</td>
<td>V</td>
<td># 3</td>
</tr>
<tr>
<td>100</td>
<td>3</td>
<td>U</td>
<td>10000</td>
<td>W</td>
<td># 4</td>
</tr>
</tbody>
</table>
The solution

<table>
<thead>
<tr>
<th>ts_quantum</th>
<th>ts_tqexp</th>
<th>ts_slpret</th>
<th>ts_maxwait</th>
<th>ts_lwait</th>
<th>LEVEL</th>
</tr>
</thead>
<tbody>
<tr>
<td>1000</td>
<td>X=0</td>
<td>1</td>
<td>50000</td>
<td>1</td>
<td># 0</td>
</tr>
<tr>
<td>500</td>
<td>Y=0</td>
<td>2</td>
<td>20000</td>
<td>2</td>
<td># 1</td>
</tr>
<tr>
<td>200</td>
<td></td>
<td>3</td>
<td>10000</td>
<td>3</td>
<td># 2</td>
</tr>
<tr>
<td>100</td>
<td></td>
<td>Z=4</td>
<td>10000</td>
<td>V=4</td>
<td># 3</td>
</tr>
<tr>
<td>100</td>
<td></td>
<td>U=4</td>
<td>10000</td>
<td>W=4</td>
<td># 4</td>
</tr>
</tbody>
</table>

- Recall that the only valid priorities are 0 to 4!
MacOS X Scheduler (I)

- Mac OS X uses a multilevel feedback queue
  - Manages threads, not processes
  - Four priority bands for threads
    - Normal
    - System high priority
    - Kernel mode only
    - Real-time
MacOS Scheduler (II)

- Thread priorities will vary
  - Must remain within their bands
  - Real-time threads tell the scheduler the number $A$ of clock cycles they will need out of the next $B$ clock cycles
    - Say 4000 out of the next 9000 clock cycles
Windows Scheduler

- An update of the old VMS scheduler
- Scheduler manages *threads* rather than processes.
- Has 32 priority levels:
  - 16 to 31 for *real-time threads*
  - 0 to 15 for *other threads*
- *Priority zero reserved* for the system thread zeroing free pages
Priority classes

- Apply to processes
- Five classes of process priorities
  - IDLE_PRIORITY_CLASS
  - BELOW_NORMAL_PRIORITY_CLASS
  - NORMAL_PRIORITY_CLASS
  - ABOVE_NORMAL_PRIORITY_CLASS
  - HIGH_PRIORITY_CLASS
  - REALTIME_PRIORITY_CLASS
Base priorities

- Apply to threads
- Defined within each process class
  - THREAD_PRIORITY_IDLE
  - THREAD_PRIORITY_LOWEST
  - THREAD_PRIORITY BELOW NORMAL
  - THREAD_PRIORITY NORMAL
  - THREAD_PRIORITY ABOVE NORMAL
  - THREAD_PRIORITY HIGHEST
  - THREAD_PRIORITY TIME_CRITICAL
Real-time threads

- Real-time processes belong to REALTIME_PRIORITY_CLASS
- Threads at fixed priorities between 16 and 31
  - Specified by their base priority
- Scheduling is round-robin within each priority level
Other threads (I)

- Run at *variable priorities* between 1 and 15
- Each thread has a *base priority*
  - Value depends on process class and thread priority level within class
    - 1 for all threads with `THREAD_PRIORITY_IDLE`
    - 15 for all threads with `THREAD_PRIORITY_TIME_CRITICAL`
Other threads (II)

- Thread priorities *never go below* their base priority
- These priorities are
  - *"Boosted"* whenever they return from the blocked state
  - *Decremented* when they exhaust their time slice
Thread affinity

- **Thread affinity** specifies the set of processors on which the thread can run.
  - "Setting thread affinity should generally be avoided because it can interfere with the scheduler's ability to schedule threads effectively across processors."

Thread ideal processor

- Instructs the scheduler to run the thread on that processor whenever possible
  - Does *not* guarantee that processor will always be chosen
Note

- Do not be confused by the two different usages of "suspended"
  - Suspending a process is the same as swapping it out
  - Suspending a thread in this context means moving it to the blocked state
Guaranteed scheduling

- Class of scheduling algorithms that want to ensure that its process has its *fair share* of CPU time

- Penalize processes that have already used a large amount of CPU

- Most versions of UNIX, Windows NT and after, Linux
Old UNIX Scheduler (I)

- Priorities take into account *past CPU usage*

\[ p_{\text{usrpri}} = \text{PUSER} + \frac{p_{\text{cpu}}}{2} + p_{\text{nice}} \]

where
- **PUSER** is the user's base priority
- **p_cpu** its current CPU usage
- **p_nice** a user-settable parameter
Old UNIX Scheduler (II)

- `p_cpu` is updated every second according to a decay function
  \[
  \text{decay}(p_{cpu}) = \frac{p_{cpu}}{2}
  \]
- After \( k \) seconds, penalty is decreased by a factor \( 1/2^k \)
BSD scheduler (I)

- The time quantum is 100 ms
  \[ p_{usrpri} = PUSER + \frac{p_{cpu}}{4} + 2 \times p_{nice} \]

- \( p_{cpu} \) is updated every second according to:
  \[ p_{cpu} = \frac{(2 \times ld)}{(2 \times ld + 1)} \times p_{cpu} + p_{nice} \]

- where \( ld \) is a sampled average of the length of the run queue over the last minute

DO NOT MEMORIZE THIS
BSD scheduler (II)

- Unlike the old UNIX scheduler, the BSD scheduler takes into account the system load.
  - Old CPU usage is forgiven more slowly when system load is high.
Linux 2.4 scheduler (I)

- Partitions the CPU time into *epochs*.
- At the beginning of each epoch, each process is assigned a *time quantum*
  - Specifies the maximum CPU time the process can have during that epoch.
- Processes that exhaust their time quantum cannot get CPU time until the next epoch starts.
Linux 2.4 scheduler (II)

- Processes that release the CPU before their time quantum is exhausted can get more CPU time during the same epoch.
- Epoch ends when all ready processes have exhausted their time quanta.
- Priority of a process is the sum of its base priority plus the amount of CPU time left to the process before its quantum expires.
Stride scheduling (I)

- Deterministic fair-share scheduler
- Each job has a stride
  - *Inversely* proportional to the number $n$ of tickets it has
  - If thread A has 10 tickets, thread B has 5 tickets and thread C has 20 tickets
    - Stride of A is 10, stride of B is 20 and stride of C is 5
Stride scheduling (II)

- Each process has a **pass** value
  - Initially set to process stride
- Each time a process releases the CPU
  - Scheduler selects process with *lowest pass*
  - Gives it the CPU for a *fixed time slide*
- Each time a process gets the CPU
  - Scheduler *adds* the process stride to its pass value
The key idea

- Still use epochs/Have a process priority ("pass")
  - Initially set to "stride"
    - Inversely proportional to the number of tickets of lottery scheduling
- Always schedule process with lowest pass
- Penalize differently
Stride scheduling (III)

- Scheme is *starvation free*
  - Processes that do not get any CPU time keep their original pass values
  - Other processes will see their pass values increase
Example

<table>
<thead>
<tr>
<th>Pass values</th>
<th>Scheduler will pick thread</th>
</tr>
</thead>
<tbody>
<tr>
<td>Thread A</td>
<td></td>
</tr>
<tr>
<td>10 tickets</td>
<td></td>
</tr>
<tr>
<td>stride is 10</td>
<td></td>
</tr>
<tr>
<td>10</td>
<td>C</td>
</tr>
<tr>
<td>10</td>
<td>C</td>
</tr>
<tr>
<td>10</td>
<td>C</td>
</tr>
<tr>
<td>10</td>
<td>A</td>
</tr>
<tr>
<td>20</td>
<td>C</td>
</tr>
<tr>
<td>20</td>
<td>C</td>
</tr>
<tr>
<td>20</td>
<td>C</td>
</tr>
<tr>
<td>Thread B</td>
<td></td>
</tr>
<tr>
<td>5 tickets</td>
<td></td>
</tr>
<tr>
<td>stride is 20</td>
<td></td>
</tr>
<tr>
<td>20</td>
<td></td>
</tr>
<tr>
<td>20</td>
<td></td>
</tr>
<tr>
<td>20</td>
<td></td>
</tr>
<tr>
<td>10</td>
<td></td>
</tr>
<tr>
<td>10</td>
<td></td>
</tr>
<tr>
<td>Thread C</td>
<td></td>
</tr>
<tr>
<td>25 tickets</td>
<td></td>
</tr>
<tr>
<td>stride is 4</td>
<td></td>
</tr>
<tr>
<td>4</td>
<td></td>
</tr>
<tr>
<td>8</td>
<td></td>
</tr>
<tr>
<td>12</td>
<td></td>
</tr>
<tr>
<td>12</td>
<td></td>
</tr>
<tr>
<td>16</td>
<td></td>
</tr>
<tr>
<td>12</td>
<td></td>
</tr>
</tbody>
</table>

10 20 4

20 20 12

20 20 16
Explanations

- Process C gets first slot
  - Lowest pass value (4)
- Process C gets second slot
  - Lowest pass value (8)
- Process A gets third slot
  - Lowest pass value (10)
- Process C gets fourth slot
  - Lowest pass value (12)
Note

- Whenever two threads have the same pass value, the scheduler will pick the thread with the *lowest stride*
FreeBSD 5.0 ULE scheduler

- Designed for threads running on multicore architectures

- Two parts
  - Low-level scheduler
    - Run every time a core is released
  - High-level scheduler
    - Run every second
Low-level scheduler

- Kernel maintains a set of *run queues* for each CPU
  - With different priorities

- Low-level scheduler selects first thread on highest-level non-empty run queue
High-level scheduler

- Reevaluates thread priorities
  - Real-time threads have fixed priorities
  - Scheduler detects interactive threads on the base of their interactivity score:
    - Scaling factor $\times \frac{\text{Sleep time}}{\text{Run time}}$
- Also assigns threads to CPUs
  - Complex process
Observations

- The low-level scheduler is kept simple
  - Makes quick decisions

- The high-level scheduler uses a very clever method to detect interactive processes

\[
\text{(Voluntary) Sleep time} \quad \frac{\text{Run time}}{\text{Sleep time}}
\]

- Must still decide length of observation period
  - Interactive processes can go through bursts of high CPU activity