Solutions to the first midterm

COSC 4330/6310 Summer 2013

a) Give an example of a popular operating system using a UNIX or a Linux kernel.

b) Give an example of a *real-time process* with *hard deadlines*.

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 Industrial process control, missile guidance system,

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 - They are slower than other kernel organizations
- d) What is the main advantage of *delayed writes*?
 - They reduce the number of disk accesses
 - They return faster

e) Where was the first operating system with a graphical user interface developed?

f) Which event(s) will move a process from the waiting state to the ready state?

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- f) Which event(s) will move a process from the waiting state to the ready state?
 - Whenever a system request issued by the process completes

What would be the *main disadvantage* of a processor that would not have (3×5 points)
 a) Separate *supervisor* and *user modes*?

b) Memory protection?

c) Timer interrupts?

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What would be the *main disadvantage* of a processor that would not have (3×5 points) a) Separate *supervisor* and *user modes*? We could not prevent user programs from directly accessing the disk **b)** Memory protection? We could not prevent user programs from corrupting the kernel c) Timer interrupts?

- What would be the *main disadvantage* of a processor that would not have (3×5 points)
 a) Separate *supervisor* and *user modes*?
 - We could not prevent user programs from directly accessing the disk
 - **b)** Memory protection?
 - We could not prevent user programs from corrupting the kernel
 - c) Timer interrupts?
 - We could not prevent user programs from monopolizing the CPU

Third question

 Add the missing code to the following program to ensure it will *always* print:
 Child says hello! Parent says hello!

in that exact order. (3×5 points)

```
#include <unistd.h>
#include <stdio.h>
void main() {
    int pid;
    if ((pid = fork()) == ___) {
        printf("Child says hello!\n");
    }
}
```

}// if

printf("Parent says hello!\n");
} // main

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#include <stdio.h>
void main() {
 int pid;
 if ((pid = fork()) == ____) {
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 }
}

} // if

printf("Parent says hello!\n");
} // main

#include <unistd.h>
#include <stdio.h>
void main() {
 int pid;
 if ((pid = fork()) == 0) {
 printf("Child says hello!\n");
 _exit(0); // to terminate the child
 } // if

printf("Parent says hello!\n");
} // main

#include <unistd.h> #include <stdio.h> void main() { int pid; if ((pid = fork()) = () { printf("Child says hello!\n"); exit(0); // to terminate the child } // if wait(0); // wait first for child completion printf("Parent says hello!\n"); } // main

What does a program do when it receives a signal? (5 points)

What can we do to change this behavior?
 (5 points)

Is it always possible? (5 points)

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What does a process do when it receives a signal? (5 points)

- It terminates
- What can we do to change this behavior?
 (5 points)
 - Process can catch the signal using signal(...) systeem call
- Is it always possible? (5 points)
 - SIGKIL signal cannot be caught

Fifth question

How will the following code fragment affect stdin, stdout and stderr? (3×5 points)

> int fda, fdb; fda = open("alpha", O_RDWR | O_CREAT, 0640); fdb = open("beta" , O_RDWR | O_CREAT, 0640); close(0); dup(fda); close(1); dup(fdb);

> int fda, fdb; fda = open("alpha", O_RDWR | O_CREAT, 0640); fdb = open("beta" , O_RDWR | O_CREAT, 0640); close(0); dup(fda); // fda duplicated into stdin close(1); dup(fdb);

Stdin will read from file "alpha"

> int fda, fdb; fda = open("alpha", O_RDWR | O_CREAT, 0640); fdb = open("beta", O_RDWR | O_CREAT, 0640); close(0); dup(fda); // fda duplicated into stdin close(1); dup(fdb); // fdb duplicated into stdout >stdin will read from file "alpha" >stdout will be redirected to file "beta"

> int fda, fdb; fda = open("alpha", O_RDWR | O_CREAT, 0640); fdb = open("beta", O_RDWR | O_CREAT, 0640); close(0); dup(fda); // fda duplicated into stdin close(1); dup(fdb); // fdb duplicated into stdout >stdin will read from file "alpha" >stdout will be redirected to file "beta" Stderr will be unchanged

Sixth question

What happens when a *user-level thread* does a *blocking system call*? (5 points)

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Sixth question

What happens when a *user-level thread* does a *blocking system call*? (5 points)

 The CPU scheduler will put the whole process into the waiting state even when other threads are ready to run

What can we do to avoid that? (5 points)

- The programmer cannot use blocking system calls
- Use kernel-supported threads