

NAME: \_\_\_\_\_ (FIRST NAME FIRST) SCORE: \_\_\_\_\_

**COSC 4330/6310**

**SECOND MIDTERM**

**APRIL 4, 2011**

*This exam is **closed book**. You can have **one page** of notes. UH expels cheaters.*

1. *Questions with short answers:* (4×5 points)

a) What would be the major disadvantage of using *non-blocking receives* to build a server?

The server will waste cycles doing busy waits between requests.

b) What is the easiest way to implement the *at most once semantics* in remote procedure calls?

Attach a serial number to each request and instruct the server to reject duplicates of previous requests.

c) What is the difference between *virtual circuits* and *streams*?

Virtual circuits preserve message boundaries; streams do not..

d) When are *busy waits* the best choice?

In multicore/multiprocessor architectures when a process waits for a process running on another processor and the wait is expected to be short.

2. Consider the instruction **TSET R7, LOCK** and assume it is used to ensure mutual exclusion within a critical section. Assuming that the variable **LOCK** can only be equal to zero or one, what are the two possible values for **R7** after the instruction is executed and their meanings? (2×5 points)

a) If R7 equals   0   then the process can enter the critical section.

b) If R7 equals   1   then the process cannot enter the critical section.

3. For each of the statements below, indicate in one sentence whether the statement is true or false (2 points), **and why** (3 points).

- a) Making all remote procedures *idempotent* greatly simplifies the task of the RPC server.

TRUE, we do not have to worry about multiple executions of the same procedure call.

- b) A *blocking send* is the same as a *buffered send*.

FALSE, a buffered send is the same as a non-blocking send.

- c) The *all or nothing semantics* guarantees that all remote procedure calls will be executed *at least once*.

FALSE, the all or nothing semantics guarantees that all remote procedure calls will be executed exactly once or not at all.

- d) We can simulate a *blocking receive* with a *non-blocking receive* inside a busy wait loop.

TRUE, think of `while (non_blocking_read(...) == NO_MESSAGE);`

- e) Peterson's algorithm *requires busy waits*.

TRUE, it contains an empty while loop.

- f) In a RPC, one of the tasks of the *user stub* is to exchange messages with the *user program*.

FALSE, user stubs exchange messages with the server stub.

4. Complete the following template to obtain a correct solution to the mutual exclusion problem for two processes whose ID's are either 0 or 1? (5×4 points)

```
shared int requested[2] = {0, 0};

shared int turn;

void enter_region(int mypid) {
    requested[____mypid____] = ____1____;
    turn = ____1 - mypid____;
    while (requested[____1 - mypid____] && turn != mypid);
} // enter_region

void leave_region(int mypid) {
    requested[____mypid____] = 0;
} // leave_region
```

5. Give at least one example of distributed applications

- a) That should use *streams* rather than *datagrams*. (5 points)

Applications requesting transfers of large amounts of data: http, ftp, ...

- b) That should use *datagrams* rather than *streams*. (5 points)

Applications requesting transfers of small amounts of data

6. Two concurrent processes access the same shared variable **count**.

```
process one {                process two {
    count++;                  count--;
}
```

Assuming that **count** was initially equal to 3, what values can it take after the two processes have completed? (10 points minus 5 points for each incorrect or missing answer)

**Answers:** 2, 3 and 4