



SOLUTIONS FOR THE SECOND 3360/6310 QUIZ

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First question

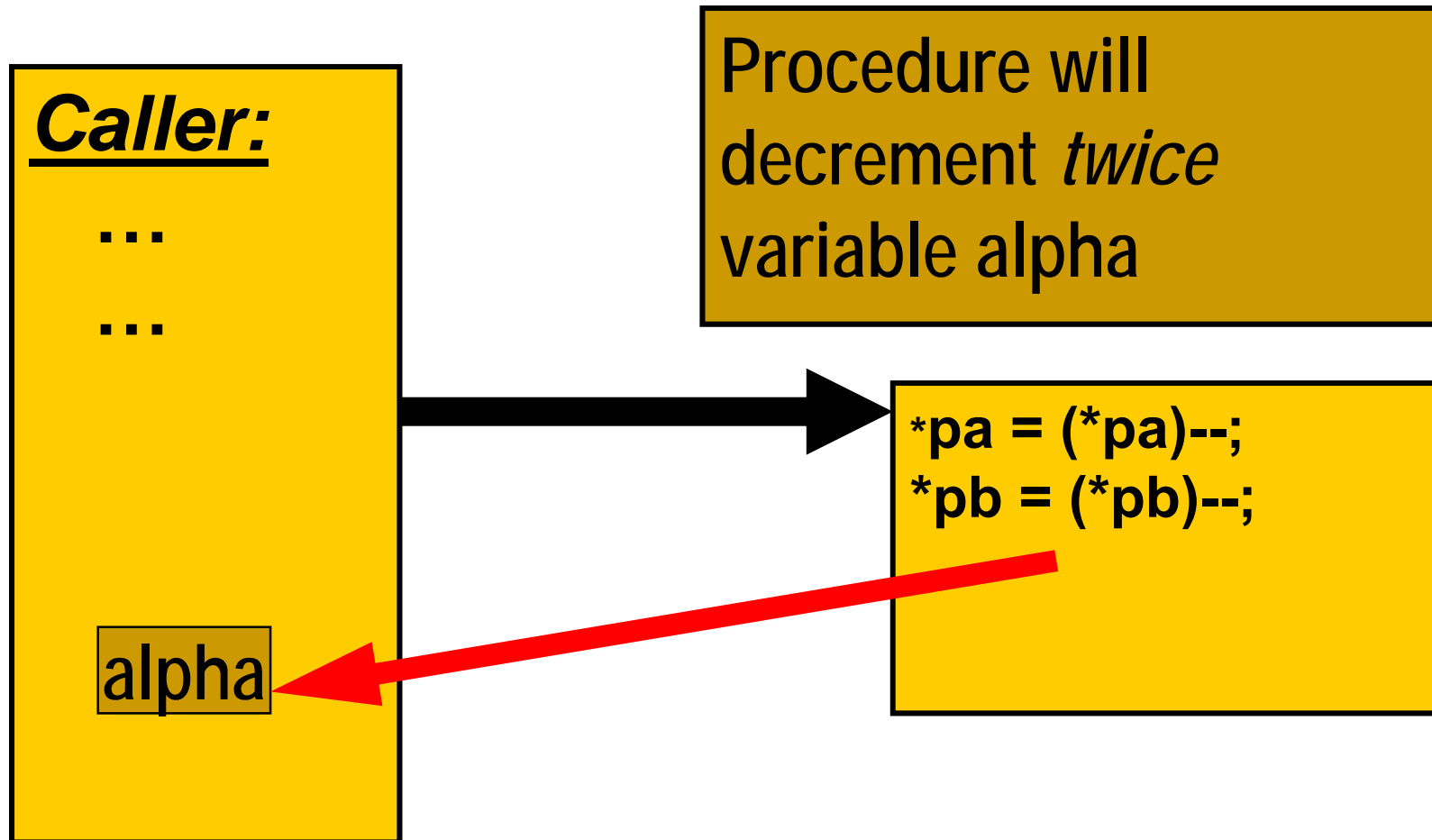
- Consider the function

```
□ void doubledecrement(int *pa, int *pb){  
    *pa = (*pa)--;  
    *pb = (*pb)--;  
} // doubledecrement
```

- and assume the following calling sequence:

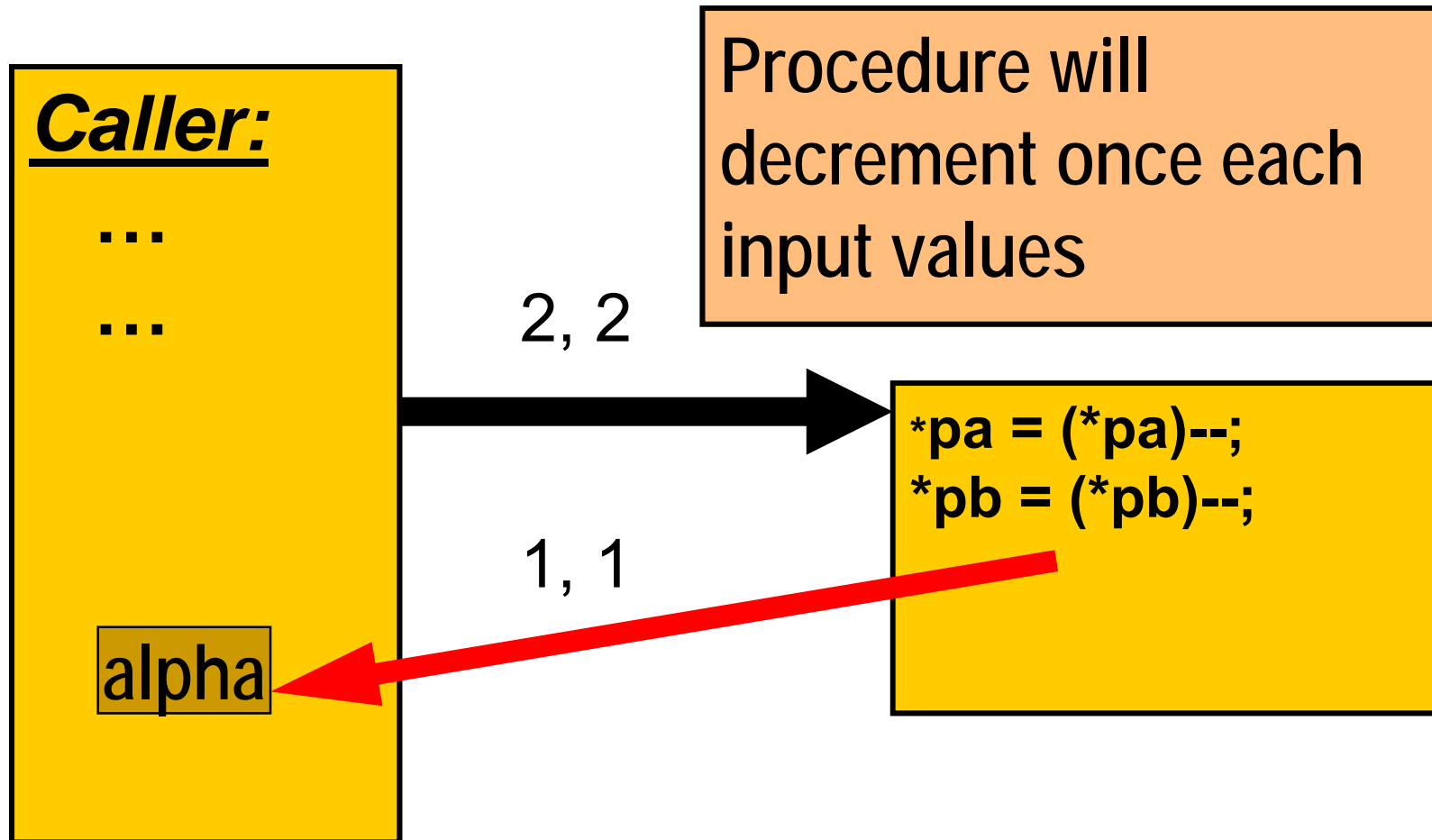
```
□ alpha = 2;  
  doubledecrement (&alpha, &alpha);
```

Passing by reference





Passing by value and result





First question

- What will be the value of **alpha** *after the call* assuming that the call was:
 - A *conventional procedure call*?
 - **alpha = 0**
 - A *remote procedure call*?
 - **alpha = 1**



Second question

- Consider the following system V.4 scheduler:

#ts_quantum	ts_tqexp	ts_slpret	ts_maxwait	ts_lwait	LVL
800	X	1	16000	0	# 0
400	0	2	8000	2	# 1
200	1	3	4000	3	# 2
100	2	3	2000	3	# 3

- Which level corresponds to the highest process priority?



Second question

- Consider the following system V.4 scheduler:

#ts_quantum	ts_tqexp	ts_slpret	ts_maxwait	ts_lwait	LVL
800	X	1	16000	0	# 0
400	0	2	8000	2	# 1
200	1	3	4000	3	# 2
100	2	3	2000	3	# 3

- Which level corresponds to the highest process priority?

□ **Level 3**



Second question

- Consider the following system V.4 scheduler:

#ts_quantum	ts_tqexp	ts_slpret	ts_maxwait	ts_lwait	LVL	
800	X	1	16000	0	#	0
400	0	2	8000	2	#	1
200	1	3	4000	3	#	2
100	2	3	2000	3	#	3

- What brings a process in the **ts_tqexp** column?



Second question

- Consider the following system V.4 scheduler:

#ts_quantum	ts_tqexp	ts_slpret	ts_maxwait	ts_lwait	LVL
800	X	1	16000	0	# 0
400	0	2	8000	2	# 1
200	1	3	4000	3	# 2
100	2	3	2000	3	# 3

- What brings a process in the **ts_tqexp** column?
 - The process returned to the ready queue after having exhausted its time slice.



Second question

- Consider the following system V.4 scheduler:

#ts_quantum	ts_tqexp	ts_slpret	ts_maxwait	ts_lwait	LVL
800	X	1	16000	0	# 0
400	0	2	8000	2	# 1
200	1	3	4000	3	# 2
100	2	3	2000	3	# 3

- Which value in the table is ***not correct***?



Second question

- Consider the following system V.4 scheduler:

#ts_quantum	ts_tqexp	ts_slpret	ts_maxwait	ts_lwait	LVL
800	X	1	16000	0	# 0
400	0	2	8000	2	# 1
200	1	3	4000	3	# 2
100	2	3	2000	3	# 3

- Which value in the table is *not correct*?
 - The value for ts_lwait at level 0



Second question

- Consider the following system V.4 scheduler:

#ts_quantum	ts_tqexp	ts_slpret	ts_maxwait	ts_lwait	LVL	
800	X	1	16000	0	#	0
400	0	2	8000	2	#	1
200	1	3	4000	3	#	2
100	2	3	2000	3	#	3

- Why?



Second question

- Consider the following system V.4 scheduler:

#ts_quantum	ts_tqexp	ts_slpret	ts_maxwait	ts_lwait	LVL
800	X	1	16000	0	# 0
400	0	2	8000	2	# 1
200	1	3	4000	3	# 2
100	2	3	2000	3	# 3

- Why?
 - It does not prevent starvation for level 0 processes



Third question

- What is the ***main drawback*** of ***virtual circuits and streams***?



Third question

- What is the ***main drawback*** of ***virtual circuits and streams***?
 - **The high setup cost of the connection.**



Third question

- What is the ***sole disadvantage*** of ***atomic transactions***?



Third question

- What is the ***sole disadvantage*** of ***atomic transactions***?
 - **Their high cost.**



Third question

- How can we prevent deadlocks by ***denying*** the ***circular wait*** condition?



Third question

- How can we prevent deadlocks by ***denying*** the ***circular wait*** condition?
 - **By requiring them to acquire all their resources in the same linear order.**



Third question

- What is the difference between ***blocking sends*** and ***non-blocking sends***?



Third question

- What is the difference between ***blocking sends*** and ***non-blocking sends***?
 - A blocking send waits until its message has been delivered to its recipient.
 - A non-blocking send does not.



Third question

- What is the ***main disadvantage*** of ***lottery scheduling***?



Third question

- What is the *main disadvantage* of *lottery scheduling*?
 - It does not guarantee that processes that received a lot of tickets will always execute ahead of the others.



Third question

- What is the ***safest place*** to put your ***signal()*** call in your monitor procedures?



Third question

- What is the ***safest place*** to put your ***signal()*** call in your monitor procedures?
 - **At the end of these procedures.**



Fourth question

- A small parking lot has space for 30 cars and a single entry/exit point that can only accommodate one car at a time.
- Complete the following solution in a way that avoids deadlocks.



Declarations

```
■ semaphore spaces = _____;  
    semaphore green = _____;
```



leave_lot() function

```
■ leave_lot(){  
    P(&green);  
    get_out();  
    V(&green);  
    V(&spaces);  
} //leave lot
```



enter_lot() function

```
■ enter_lot(){  
    _____;  
    _____;  
    get_in();  
    _____;  
} // enter_lot
```



Declarations

■ semaphore spaces = 30;

semaphore green = 1; // our mutex



enter_lot() function

```
■ enter_lot(){  
    P(&spaces);  
    _____;  
    get_in();  
    _____;  
} // enter_lot
```



enter_lot() function

```
■ enter_lot(){  
    P(&spaces);  
    P(&green);  
    get_in();  
    V(&green);  
} // enter_lot
```




Fifth question

- What is the major disadvantage of *busy waits*?
- What can we do to eliminate them?
- Are there cases where it is better to keep them?



Answers

- ☐ **Busy waits waste CPU cycles and cause unnecessary context switches.**
- ☐ **We should use instead kernel-supported solutions that move the waiting processes to the blocked state.**
- ☐ **For short waits in multicore architectures.**