



# **SOLUTIONS FOR THE THIRD 3360/6310 QUIZ**

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# First question

- What is the main advantage and the main disadvantage of ***access control lists*** over ***tickets***?
  - ***Main advantage of access control lists:***
  - ***Main disadvantage of access control lists:***



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- What is the main advantage and the main disadvantage of ***access control lists*** over ***tickets***?
  - ***Main advantage of access control lists:***
    - ***They are more flexible than tickets***
  - ***Main disadvantage of access control lists:***
    - ***They are much slower***



## Second question

- What does an FFS *cylinder group* contain?
- How do these cylinder groups improve the performance of the file system?



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  - *They contain both i-node blocks and data blocks*
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  - *They contain both i-node blocks and data blocks*
- How do these cylinder groups improve the performance of the file system?
  - *They eliminate most long seeks by allocating to files data blocks in the same cylinder group as their i-nodes.*



## Third question

- A virtual memory system has **48-bit** addresses.  
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    - **$2^{13} = 8,192$  bytes**



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Given that **13** of these **48** bits are used by the bit offset,
- How many bits of the virtual address are used by the page number?
  - **$48 - 13 = 35$  bits**



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  - On average, how much memory is lost to internal fragmentation?
- *Half a page or 4,096 bytes*



## Fourth question

- What is the main purpose of the UNIX **mount** primitive?



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- What is the main purpose of the UNIX **mount** primitive?
  - *To make the directory tree of a given disk partition appear as a subtree of the directory tree of another partition.*
  - *To merge the directory trees of multiple disk partitions into a single tree.*



## Fourth question

- What is the major disadvantage of using ***large blocks*** in a file system?





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- What is the major disadvantage of using ***large blocks*** in a file system?
  - ***They increase internal fragmentation.***



## Fourth question

- What is the major disadvantage of letting the kernel handle ***TLB misses***?



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- What is the major disadvantage of letting the kernel handle ***TLB misses***?
  - ***Each TLB miss will require two context switches.***



## Fourth question

- What is the main advantage of ***journaling file systems*** using ***asynchronous log writes***?



## Fourth question

- What is the main advantage of ***journaling file systems*** using ***asynchronous log writes***?
  - ***They are much faster than journaling file systems using synchronous log writes.***



## Fourth question

- Why did BSD FFS introduce ***block fragments***?



## Fourth question

- Why did BSD FFS introduce ***block fragments***?
  - ***To reduce internal fragmentation.***



## Fourth question

- When a user creates a ***new file***, what should be written to disk ***first*** (a) the ***directory block*** containing the new entry for the file ***or*** (b) the ***i-node block*** containing the new i-node?





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- When a user creates a ***new file***, what should be written to disk ***first*** (a) the ***directory block*** containing the new entry for the file ***or*** (b) the ***i-node block*** containing the new i-node?
- ***We should write first the i-node block containing the new i-node.***



# Fifth question

- A virtual memory system uses **4KB** pages and ***inverted page tables***. Assuming that each page table entry occupies **24 bytes** what is the fraction of main memory occupied by page tables?

## ***Why 24 bytes?***

*Assuming 64-bit addresses, we will have 8 bytes for the page number, 8 bytes for the page frame number and 8 bytes for the address of the the next PTE in the bucket.*



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□  $24/4,096 = 3/512 = 0.586 \text{ percent}$



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  - What happens when the *first hand* of the clock reaches a *valid page*?



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- Consider the *two-handed BSD clock replacement policy*.
  - What happens when the *first hand* of the clock reaches a *valid page*?
    - *It marks it invalid.*



## Sixth question

- Consider the *two-handed BSD clock replacement policy*.
  - What happens when the *second hand* of the clock reaches a *valid page*?



## Sixth question

- Consider the ***two-handed BSD clock replacement policy***.
  - What happens when the ***second hand*** of the clock reaches a ***valid page***?
    - ***Absolutely nothing.***



## Sixth question

- Consider the *two-handed BSD clock replacement policy*.
  - What happens when the *second hand* of the clock reaches a page that was *marked invalid*?





## Sixth question

- Consider the *two-handed BSD clock replacement policy*.
  - What happens when the *second hand* of the clock reaches a page that was *marked invalid*?
    - *It expels it.*



# Seventh question

- A **32-bit** Berkeley UNIX file system has a block size of **eight** kilobytes.
- How many indirect blocks will be accessed through one level of indirection:
  - ☐ For a **54-kilobyte** file?
  - ☐ For a **two-megabyte** file?



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  - ☐ For a **54-kilobyte** file? ***None***
    - ***(Can access 96 KB straight from i-node)***
  - ☐ For a **two-megabyte** file? ***244***
    - ***File has 256 data blocks***
    - ***The first 12 are accessed straight from the i-node***