Chapter VI DEADLOCKS (short version)

> Jehan-François Pâris jfparis@uh.edu

Overview

- Deadlocks
- Necessary conditions for deadlocks
- Deadlock prevention

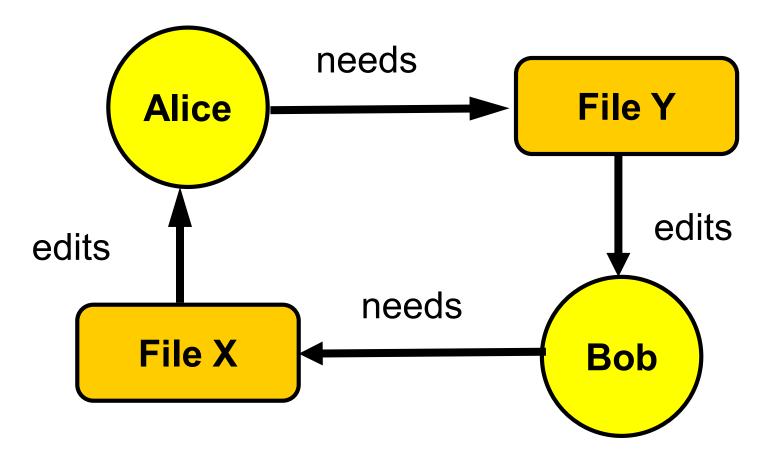
Deadlocks

 A deadlock is said to occur whenever
 Two or more processes are blocked
 Each of these processes is waiting for a resource that is held by another blocked process.

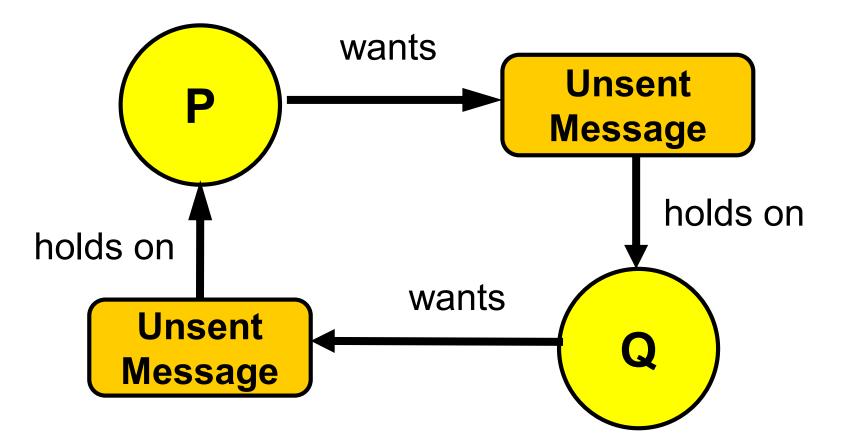
Examples

- Alice edits file X and needs to access file Y
 Bob edits file Y and needs to access file X
- Process P expects a message from process Q
 Process Q expects a message from process P

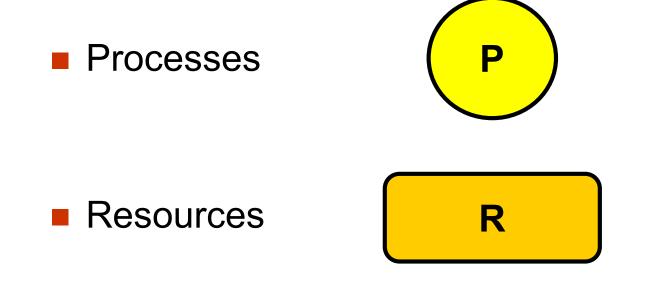
A graphic view (I)



A graphic view (II)

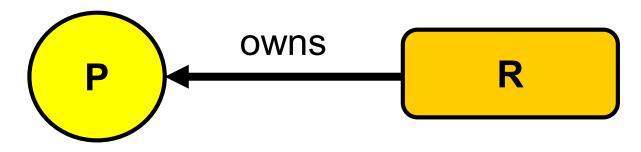


Elements

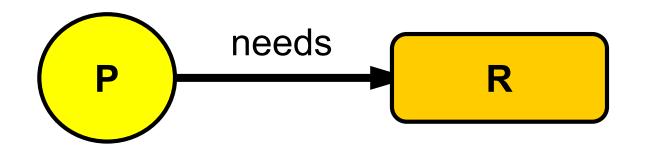


Relations

Process P holds on/owns resource R



Process P needs/wants resource R



Serially reusable resources

- Memory space, buffer space, disk space, USB slot to insert a flash drive
- Exist only in a *limited quantity*
- One process may have to *wait* for another process to release the resources it needs.

Consumable resources

- Cannot be reused
- Messages are best example:
 - "Owned" by the process that creates them until it releases them
 - □ "Wanted" by the process that waits for them

Handling deadlocks

- Do nothing: Ignore the problem
- Deadlock prevention: Build deadlock-free systems



- Deadlock avoidance: Avoid system states that could lead to a deadlock
- Deadlock detection: Detect and break deadlocks

Haberman's conditions

- Four necessary conditions must all be in effect for deadlocks to happen:
 - Mutual Exclusion
 Hold and Wait
 No Preemption
 Circular Wait

Mutual exclusion

At least one of the processes involved in the deadlock must claim exclusive control of some of the resources it requires

□ No sharing

Hold and wait

Processes can hold the resources that have already been allocated to them while waiting for additional resources

No preemption

Once a resource has been allocated to a process, it cannot be taken away or borrowed from that process until the process is finished with it

Circular wait

- There must be a circular chain of processes such that each process in the chain holds some resources that are needed by the next process in the chain.
 - □ Formal equivalent to what we call a **vicious circle**

Deadlock prevention

- Any system that prevents any of the four necessary conditions for deadlocks will be deadlock-free
- Must find the easiest condition to deny

Denying mutual exclusion

Prevent any process from claiming exclusive control of any the resource

Drawbacks

Many resources can only be used by one process at a time
 Cannot hold on a message and send it at the same time

Denying hold and wait

Require processes to get *all* the resources they will need or *none of them*

Drawbacks

- Forces processes to acquire ahead of time all the resources they might need
- □ Does not apply the consumable resources such a messages

Allowing preemption

Let processes take away or borrow the resources they need from the processes that hold on them

Drawbacks

□ Will result in *lost work* when a process steals storage space from another process

□ Cannot force processes to send messages

Denying circular wait (I)

- Impose a total order on all resource types and force all processes to follow that order when they acquire new resources
- If a process needs more than one unit of a given resource type it should acquire all of them or none

Denying circular wait (II)

Works very well for resources like CPU and memory

Drawbacks

□ Would force messages to move in only one direction

Processes could not exchange messages

A little problem

- Two courses at South Hillcroft University are *co-requisites* of each other:
 - □ Sandcastle Design
 - □ History of Sandcastles
- Both courses are likely to be oversubscribed
- Think of possible deadlocks
- Post your solutions on Prulu

Check list

- You must understand
 - Difference between consumable and serially-reusable resources
 - □ Haberman's four necessary conditions for deadlocks
 - □ The four ways to deny them
 - Why they do not work for client/server systems