Na	ame:			_ (First name first	i) S	core:	
C	DSC 6360		Quiz #1		SEPTEMB	ER 14 ,	2011
Clo	osed book. You c	an have with you on	e single-sided 8	3½ by 11 sheet of r	notes. UH ex	kpels ch	eaters
1.	How many lines wi	II be printed by the follo	wing program? (5	points)			
	main() {	fork(); fork()	; printf("Do	one!\n") }			
	Answer:	4	lines				
2.	A Berkeley UNIX fil accessed :	le system has a block s	ize of 16 kilobytes	. How many blocks	of a given file	can be	
	a) using the block	addresses stored in th	e i-node? (5 points	s)	12	blocks	
	b) with one level of	of indirection? (5 points))	16 KB/4 = 4K=	4,096	blocks	
	c) with two levels	of indirection? (5 points	s) <u>4GB/16</u>	6KB - 4,096 -12 = 2	56K- 4,108	blocks	
3.	What is the main p	urpose of the UNIX <i>set</i>	user-id bit? (10 p	oints)			
		t allows an executable hose of the user exec		the user-id—and the	e privileges—	of its	
4.	Where does UNIX	store the <i>name</i> of a file	? (10 points)				
	UNIX stores the file.	name(s) of a file in the	e directory entry	entries pointing to	the i-node o	f the	
5.		and Joy decide to limit e pages to 300 pages p			ck sweeps the	e	
			, ,	·			
	marked invalid by overhead will never	mplementation of the the Clock hand occasi er take more than 10% the clock sweeps the	ons two context : % of the total <i>C</i> P	switches. To ensur U time, Babaoğlu a	e that t this nd Joy limite	context ed the <i>sp</i>	switch

6.		What makes the choice of a block size for a file system so difficult? (10 points) How did the Unix Fast File System (FFS) solve that issue? (10 points)				
		Most file systems contain a large number of small files as well as a significant number of large and very large files. The block sizes that would be the most effective for large files would create an excessive amount of <u>internal fragmentation</u> when they are used to store to small—and very small—files.				
		UNIX FFS requires a minimum block size of 4 KB but allows blocks to be divided into 2, 4, or 8 <u>fragments</u> that can be used to store small files and the tails of larger files.				
7.	The	e Unix Fast File System used synchronous updates for all metadata updates.				
	a)	What is the <i>main advantage</i> of this approach? (10 points)				
		guarantees that the file system will remain in a consistent state after a crash and ensures e immediate durability of metadata updates.				
	b)	What is its <i>main disadvantage</i> ? (10 points)				
	Blo	cking metadata updates result in numerous seeks that slow down the file system				
8.	Wh	nat is the main motivation for <i>cylinder groups</i> ? (10 points)				
		nce each cylinder group contains both a fragment of the i-node table, subdividing a disk partition into inder groups eliminates most long seeks between the i-node table entries and the blocks they point				