

# **Solution to the First COSC 6360 Quiz for Fall 2013**

Jehan-François Pâris  
jfparis@uh.edu



# First question

- What is the purpose of the UNIX mount system call ? (10 points)

# Answer

- The mount() system call makes a file system appear as a subdirectory of another file system.

or

- It glues together the directory trees of the system disk partitions to form a single directory hierarchy.

## Second question

- In a 32-bit UNIX file system, what is the *minimum block size* that lets users access the whole contents of a file using **two levels of indirection**? (10 points)

# Answer

- With two levels of indirection and a block size  $b$ , we can address
  - $(b/4)^2 \times b$  bytes
- We must have  $(b/4)^2 \times b \geq 2^{32}$ 
  - $b^3 / 2^4 \geq 2^{32}$
  - $b^3 \geq 2^4 \times 2^{32} = 2^{36}$
  - $b \geq 2^{36/3} = 2^{12} = \underline{4K}$

# Third Question

- Why does the Fast File System subdivide each disk partition into *cylinder groups*? (20 points)

# Answer

- It minimizes disk arm motions because each cylinder group has its own fragment of the i-node table. As a result most file blocks reside closer to the file i-node.
  - (*This is not true for large files but remains true for most file blocks in the cylinder group.*)

## Fourth question (Part I)

- What is the main advantage of *copy-on-write*? (10 points)



# Answer

- Copy-on-write reduces the cost of the UNIX `fork()` system call by letting the parent and child process share their data segment and only duplicating pages that are modified by one of them.
  - *Can also mention its use for implementing efficient message passing.*

## Fourth question (Part II)

- Can you think of a case where this advantage would not hold?  
(10 points)

# Answer

- Copy-on-write will not work as well if one of the two processes modifies most of the contents of its address space between the `fork()` and the `exec()`.

## Fifth question

- Why did Babaoğlu and Joy decide not to use the VMS page replacement policy? (2x10 points)

*A sure indication that a two-part answer is expected*

# Answer

- Estimating the right size for the policy resident set of each new process was very difficult under UNIX.
- Supporting real-time processes was not an objective (*in the paper*)

## Sixth question

- In the Mach virtual memory system, which should be the protection attributes for the address range containing the data segment? (10 points)

# Answer

- Read and write but *never* execute
  - *You were asked the access rights not the inheritance attributes!*