SOLUTIONS TO THE FIRST 6360 QUIZ

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First question

Why is the UNIX fork() system call so expensive?



- Why is the UNIX fork() system call so expensive
 - Because it has to
 - Create a new address space
 - Populate it with all the pages that were present in the parent's address space



Second question

How many lines will the following program print?

```
#include <stdio.h>
#include <unistd.h>
void main() {
    fork(); fork();
    printf("Done!\n");
}
```



How many lines will the following program print?

```
#include <stdio.h>
#include <unistd.h>
void main() {
    fork(); fork();
    printf("Done!\n");
}
```

Four lines



Alternate second question

How many lines will the following program print?

```
#include <stdio.h>
#include <unistd.h>
void main() {
    fork();
    printf("Done!\n");
}
```



How many lines will the following program print?

```
#include <stdio.h>
#include <unistd.h>
void main() {
    fork();
    printf("Done!\n");
}
```

Two lines



Third question

What is the purpose of the UNIX set-userid bit?



- What is the purpose of the UNIX set-userid bit?
 - □ It specifies that a given program must be executed
 - With the access rights of its owner
 - Instead of the access rights of who started the program



Fourth question

What is the purpose of block fragments in the Unix Fast File System?



- What is the purpose of block fragments in the Unix Fast File System?
 - □ To reduce internal fragmentation
 - Allows the file system to allocate half block and quarter blocks to small files and the tail ends of larger files.



Fifth question

Where does Linux store its file access control lists?



Where does Linux store its file access control lists?

□ In the i-node of each file



Sixth question

- A 32-bit FFS file system has a block size of 4 kilobytes.
- What is the size of the largest file that can be accessed :
 - □ Directly from the file i-node?
 - With at most one level of indirection?

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- A 32-bit FFS file system has a block size of 4 kilobytes.
- What is the size of the largest file that can be accessed :
 - □ Directly from the file i-node?
 - First 12 blocks of the file
 - 12×4KB = 48KB

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- A 32-bit FFS file system has a block size of 4 kilobytes.
- What is the size of the largest file that can be accessed :
 - With at most one level of indirection?
 - Must add (4KB/4)×4KB = 4 MB
 - Total is 4MB + 48 KB



Alternate sixth question

- A 32-bit FFS file system has a block size of 8 kilobytes.
- What is the size of the largest file that can be accessed :
 - □ Directly from the file i-node?
 - With at most one level of indirection?

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- A 32-bit FFS file system has a block size of 8 kilobytes.
- What is the size of the largest file that can be accessed :
 - Directly from the file i-node?
 - First 12 blocks of the file
 - 12×8KB = 96KB

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- A 32-bit FFS file system has a block size of 4 kilobytes.
- What is the size of the largest file that can be accessed :
 - With at most one level of indirection?
 - Must add (8KB/4)×8KB = 16 MB
 - Total is 16MB + 96 KB



Seventh question

What is the main disadvantage of using the page valid bit to simulate a missing page reference bit?



- What is the main disadvantage of using the page valid bit to simulate a missing page reference bit?
 - □ Setting the simulated page-referenced bit back to one requires kernel intervention
 - Two context switches



Seventh question

Consider the two-handed BSD clock replacement policy.

□ What happens when the *first hand* of the clock reaches a *valid page*?



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□ What happens when the *first hand* of the clock reaches a *valid page*?



Consider the two-handed BSD clock replacement policy.

□ What happens when the *first hand* of the clock reaches a *valid page*?

It marks it invalid



Consider the two-handed BSD clock replacement policy.

□ What happens when the **second hand** of the clock reaches a **valid page**?



Eighth question

Consider the two-handed BSD clock replacement policy.

□ What happens when the *first hand* of the clock reaches a *valid page*?

It ignores it.



Consider the two-handed BSD clock replacement policy.

□ What happens when the second hand of the clock reaches a page that was marked invalid?



Consider the two-handed BSD clock replacement policy.

□ What happens when the second hand of the clock reaches a page that was marked invalid?

It expels it.