## Solution to the Fall 2012 Second COSC 6360 Quiz

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# ARC

 Consider a very small cache that can hold only 4 pages and assume that the cache is managed by the ARC replacement policy

	L1	L2
In cache	4	3
	5	2
Simulated	1	6
	7	9

How would *target\_T*1 be affected if the next page to be referenced is: a) Page 5? b) Page 6?

 Consider a very small cache that can hold only 4 pages and assume that the cache is managed by the ARC replacement policy

	L1	L2
In cache	4	3
	5	2
Simulated	1	6
	7	9

How would *target\_T1* be affected if the next page to be referenced is:
a) Page 5? No change

b) Page 6? Decreased

# **Similar question**

 Consider a very small cache that can hold only 4 pages and assume that the cache is managed by the ARC replacement policy

	L1	L2
In cache	4	3
	5	2
Simulated	1	6
	7	9

How would *target\_T1* be affected if the next page to be referenced is:
 a) Page 7?

b) Page 2?

 Consider a very small cache that can hold only 4 pages and assume that the cache is managed by the ARC replacement policy

	L1	L2
In cache	4	3
	5	2
Simulated	1	6
	7	9

How would *target\_T1* be affected if the next page to be referenced is:
 a) Page 7? Increased

b) Page 2? No change

### **TLB** issues

 A computer has a TLB with 128 entries and uses 8 KB pages. What is the *coverage* of its TLB?

- A computer has a TLB with 128 entries and uses 8 KB pages. What is the coverage of its TLB?
  - **Coverage = 128×8 KB = 1 MB**
- If page size was **4KB** 
  - Coverage = 128×4 KB = 512 KB

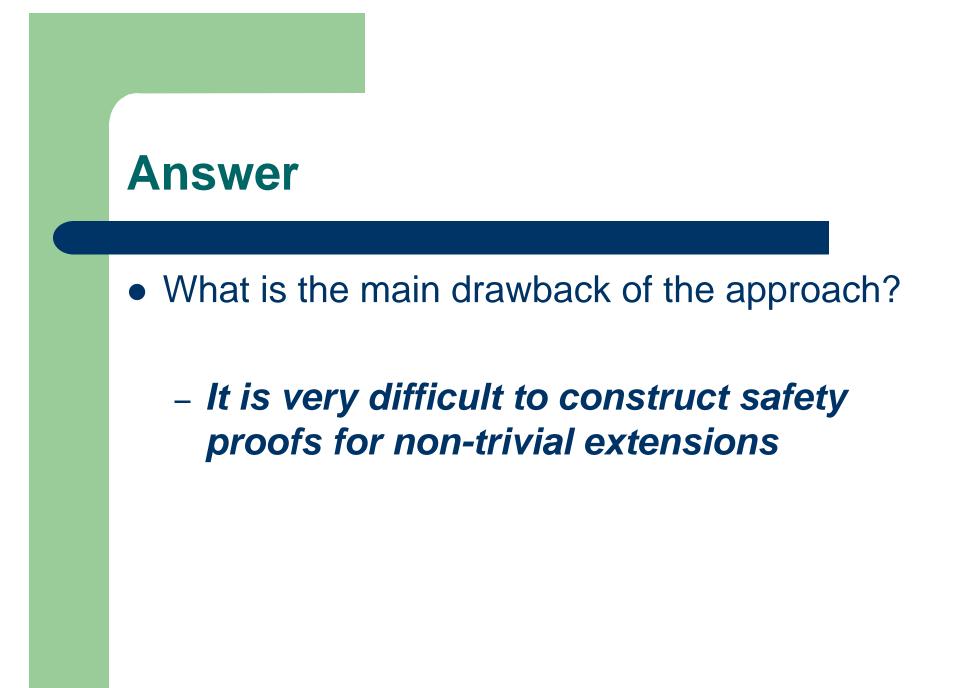
# **Proof-carrying code (I)**

• Which problem does *proof-carrying code* (PCC) attempt to solve?

- Which problem does *proof-carrying code* (PCC) attempt to solve?
  - To let kernels to check extension safety a extension load time
    - The idea is to eliminate any kind of runtime checking overhead

# **Proof carrying code (II)**

• What is the main drawback of the approach?



# Mach VM system (I)

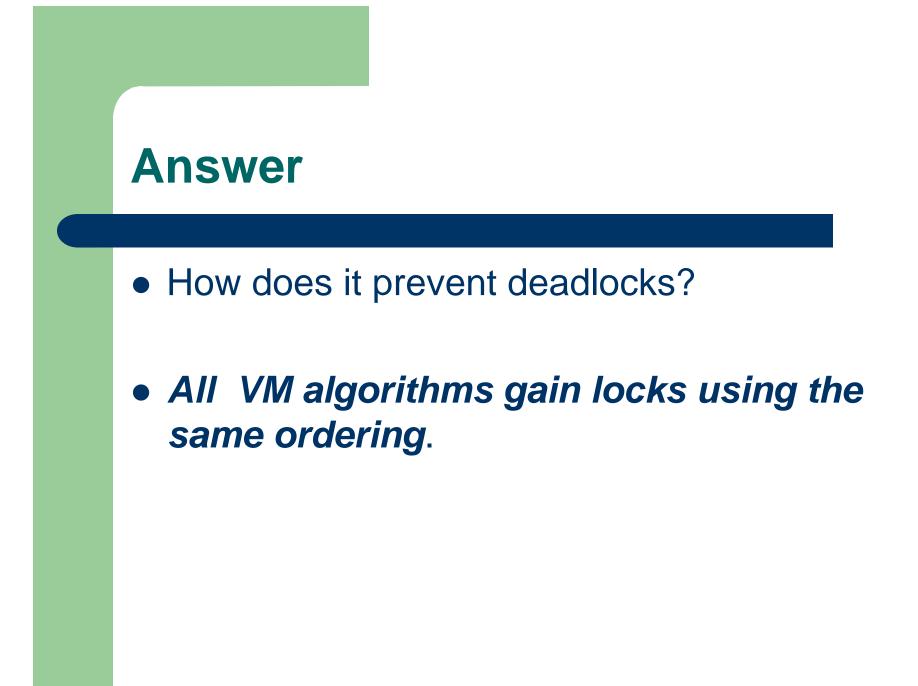
 How does the Mach virtual memory subsystem guarantee exclusive access to its data structures?

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 It uses locks whenever exclusive access to a data structure has to be guaranteed.

# Mach VM system (II)

• How does it prevent deadlocks?



## **More Mach**

• Give one example of a Mach memory object.



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- A swap area
- Any file of the file system
  - Could be
    - Executable file for a running program
    - A file being mapped

# **Treadmarks (I)**

 Treadmarks is said to use a *lazy release policy*. Why?

- Treadmarks is said to use a *lazy release policy*. Why?
  - Delays propagation of release values until an acquire is issued

# **Treadmarks (II)**

 What would be the corresponding *eager* release policy?

 What would be the corresponding *eager* release policy?

 The policy would immediately propagate all released values to all other processes

