(First name first)	
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SCORE:

COSC 6360

Quiz #3

AUGUST 2, 2010

Closed book. You can have with you one single-sided 8½ by 11 sheet of notes.

- 1. What steps must be taken by a *log-structured file system* to locate an i-node when it reboots after a *normal shutdown*? ((3×5 points)
 - a) Access <u>check-point area</u> to get address of the appropriate block of

i-node map_____

- b) Fetch that block to get address of the block containing __i-node___
- 2. Which are the main advantage and the main disadvantage of using journaling with *asynchronous log updates* compared to using journaling with *synchronous log updates*?
 - a) Main advantage: (5 points)

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Asynchronous log updates require fewer disk accesses and provide higher file system throughputs.

b) Main disadvantage: (5 points)

They do not guarantee the durability of metadata updates.

3. In Totem, what is the purpose of guaranteed vector messages? (5 points)

Guaranteed vector messages are part of the Totem multiple-ring protocol. They allow messages to be delivered to their destinations in a timely fashion without waiting for messages from a ring that remains silent.

4. In Kerberos, how does the ticket granting service (TGS)—and all other services—distinguish between valid authenticators and authenticators that are being *replayed by intruders*? (2×5 points)

TGS will reject all authenticators whose timestamps are either too old or duplicates of a timestamp ____

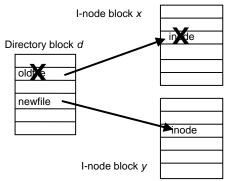
that it already received. _____

- **5.** Consider a RAID level 5 disk array with four data blocks (b_0 , b_1 , b_2 and b_3) and one parity block p per stripe. What is the most efficient way to update block b_2 and leave the array in a consistent state? (4×5 points)
 - a) Read_ values of old block b2 and old parity p. _____
 - b) Write_new value of block b₂ and new parity p = old parity p XOR nold b₂ XOR new b₂.

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6. Consider a file system using soft updates and assume that one directory block in its I/O buffer reflects the result of one file deletion and one file creation. Assuming that the i-nodes of the two files reside in two different i-node blocks, enumerate the steps the system will take to update the disk copies of the three blocks. (3×5 points)



- a) Write i-node block y.
- b) Write directory block d.
- c) Write i-node block x.
- 7. What is the main limitation of *soft updates*? (5 points)

They do not guarantee the durability of metadata updates.

- 8. Why do all NFS servers implement a *write-through policy*? (5 points) What is the *main drawback* of that policy? (5 points) Which new feature of NFS v. 3 addresses that issue? (5 points) How? (5 points)
 - a) NFS servers implement a write-through policy in order to remain stateless.
 - b) The policy results in many more disk accesses than in a regular UNIX file sytem.
 - c) NFS now has safe asynchronous writes.
 - d) Safe asynchronous writes allow NFS servers to implement delayed writes as long as the client keeps a local copy of the blocks it sends to the server. A new commit primitive lets clients check with the server that it actually has written the data.