## Solutions for Fourth Quiz

COSC 6360 Fall 2015

### First question

- Consider a diskless client trying to access a file named "/usr/joe/6360/paper.doc" that is stored on its NFS server.
  - Assuming that the client already has a handle for its root directory, how many lookup() requests will it issue?
    - Answer: \_\_\_\_\_\_ requests
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  - Assuming that the client already has a handle for its root directory, how many lookup() requests will it issue?
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  - What does NFS do to speed up these requests?
    - NFS lets client cache file handles

#### Second question

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- What does NFS do to detect them?
  - The NFS server compares the generation number of the file handle with the generation number of the i-node. If the numbers do not match, the file handle is stale.

- Consider a distributed file system implementing close-to-open consistency. Assuming that
  - Alice opens a file at 9:30 AM, modifies it and closes it at 10:15 AM,
  - Bob opens the same file at 10:00 AM, modifies it and closes it at 10:30 AM,
  - Carol opens the same file at 10:20 AM, modifies it and closes it at 11:30 AM,

Which of these three users would see his or her changes actually incorporated in the final version of the file?

#### □ Alice □ Bob □ Carol

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  - Would the server of this file system still be stateless?
    - NO, because the server will have to maintain information on leases

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  - Would that make the service *less robust*? Why?
    - No, all lease information will be obsolete by the time a failed NFS server will restart

#### Fifth question

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- Farsite encrypts user private keys with a symmetric key derived from user password and stores them in a globally-readable Farsite directory
- This frees the user from the burden of managing a fairly long non-mnemonic key.

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CRUSH is publically available because Ceph clients <u>must</u> use it to locate object replicas

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How does it simplify the design of the Ceph metadata server cluster?

Thanks to CRUSH, the Ceph metadata server does not have to maintain object lists.





Client uses CRUSH and data provided by MDS cluster to find the file

