# SOLUTIONS TO THE FIFTH 6360 QUIZ

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# First question

- Why is the Intel x86 architecture especially hard to virtualize?
  - Because its instruction set includes instructions that produce different results when executed in user mode and in privileged mode.
  - (Virtualization is easier when all privileged instructions are caught by a trap when one tries to execute them in user mode.)

# First question

How does VMWare solve this problem?

VMWare uses dynamic "binary translation" when direct execution of code would not work

# First question

How does Xen solve this problem?

Xen exports a virtual machine abstraction that is "similar but not identical to the underlying hardware"

- Paravirtualization
- Requires some modifications to the guest OS

# Second question

What does Xen do to minimize TLB flush overhead?

 The top 64MB region of each address space is reserved to Xen
Can execute Xen code without changing the page map and flushing the TLB

# Third question

Why does FAWN use a log-structured organization for its datastores?

Because it uses flash memory for its data stores and small random writes are very expensive on flash

## Fourth question

 FAWN only stores a small fragment of the key of each item in its *in-memory hash tables*.
What is the main advantage of this solution?

> It reduces the memory footprint of the hash tables thus reducing the node main memory requirements

## Fourth question

- FAWN only stores a small fragment of the key of each item in its *in-memory hash tables*.
  What is the resulting performance penalty?
  - It results in additional accesses to the secondary store
  - With the 15-bit key fragment, only one in 32,768 retrievals from the flash will require fetching an additional record.

# Fifth question

How does FAWN *delete* datalog entries?

□ It marks hash table entry invalid and adds a delete entry to the log (for durability)

# Sixth question

Why does GFS not use data caching?

Most applications stream through huge files
It would be ineffective

## Seventh question

 How does a GFS recover *after a crash*: (one-line answers)

□ Its file to chunk mappings?

By replaying its operation log

□ The locations of chunk replicas?

By polling its chunk servers