# Safely and Efficiently Multiprogramming a 64kB Computer

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June 14th, 2017

A few months before SOSP

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Emerging class of embedded applications are software platforms, rather than single purpose devices.

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## Embedded Software

- No isolation between components
- Deeply coupled components
- Static memory allocation to avoid unrecoverable runtime memory exhaustion

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Fixed concurrency at compile-time

### Embedded Hardware

Low-power budget—micro-amps average current consumption

- 64kB of RAM (to limit the sleep current)
- Memory Protection Unit—a limited hardware protection mechanism

## Challenges

Software platforms must support multiple independent dynamically-loadable apps

- -- New performance and safety concerns
- How to isolate components despite minimal hardware resouces?
- How to replace individual components without restarting the whole system?

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How to avoid fixed concurrency with limited memory

Give up on isolation—write completely bug-free code

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Whole system updates only

Give up on isolation—write completely bug-free code

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- Whole system updates only
- Use \*nix et al—forget about low power

## Tock

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**Tock** is a new operating system for low-power platforms that takes advantage of the limited hardware-protection mechanisms available on recent microcontrollers and the type-safety features of the Rust programming language to provide a multiprogramming environment:

- Isolation of software faults
- Efficient memory protection and management for dynamic application workloads
- Update/restart/remove individual (user-space) components independently
- Retains dependability requirements of long-running devices.
   Often run unattended

## Tock Architecutre



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# Capsules

- Capsules are components in the kernel
- Minimal runtime overhead:
  - Isolated "at compile-time" using the Rust language type/module system
  - Cooperatively scheduled
  - Can eliminate most isolation at compile-time

### Capsules can...

#### Most capsules are not trusted

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- Violate real-time guarantees
- Panic (sort of... lets talk...)

#### But they cannot...

- Read arbitrary memory (secret encryption keys)
- Communicate with peripherals it's not allowed to

Stronger memory isolation than hardware protection?

```
struct DMAChannel {
    ...
    enabled: bool,
    buffer: &'static [u8],
}
```

Typing hardware register can constrain allowed values with very fine granularity.

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-- Important because DMA hardware can manipulate data at any address

#### Processes

Can be unreliable since the system can respawn or kill processes without affecting other functionality.

- Hardware isolated concurrent executions of programs
- ▶ Written in any language (currently C, C++, Lua and Rust-ish)
- Total control over their memory, including dynamic heap allocation.
- Similar to processes in other systems.
  - Separate stacks allows preemptive execution
  - Memory isolated by the hardware
- Interact with kernel over a small but flexible system-call interface:
  - command, subscribe, allow
  - yield, memop

**Non-blocking API** 



command: passes a a word-size integer to a capsule allow: passes data buffers from processes to capsules subscribe:sets up a callback

What happens when the kernel requires dynamic resources to respond to a request from a process?

- We want to allow arbitrary apps so we don't know concurrency requirements:
  - How many timers will an application need?
  - Will it use SPI, UART, USB, Bluetooth, etc? One socket? 1000 sockets?
- If the kernel allocates memory for requests dynamically, it may run out of resources.

Kernel RAM	Syscall RAM	Max Used
3506	712	158
4216	1422	316
4928	2134	474
	Kernel RAM 3506 4216 4928	Kernel RAM         Syscall RAM           3506         712           4216         1422           4928         2134

TOSThreads has low memory efficiency. Static allocation costs 710-712 bytes per thread, of which at most 158 bytes (22%) can be in use at any time. These numbers do not include the thread stacks, each of which can be less than 100 bytes.

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TOSThreads, is a threads package for TinyOS that combines the ease of a threaded programming model with the efficiency of an event-based kernel. Tock allows a process to "grant" to the kernel portions of its own memory, which the kernel can use to maintain state for process requests.

- Separate sections of kernel heap located in each process's memory space.
- Grant allocations for one process do not affect kernel's ability to allocate for another.
- Type-safe interface guarantees all grants for a process can be freed immediately if the process dies.

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Basic idea: kernel API ensures there are no long-lived pointers directly to grant-allocated memory. Processes (Any language)





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- Process cannot access grant allocated memory
  - ► We use an additional, dynamically determined MPU rule
- Ensure grant-allocated values unavailable to capsules once process dies through limited API:
  - Capsules pass a closure to the enter method
  - Memory in a grant region only accessible from within closure

- Pointers to grant memory cannot escape the closure
- Implications on kernel design: should avoid cross process data-structures

```
impl<T: Default> Grant {
   fn create() -> Grant<T>
```

ł

```
fn enter<F,R>(&self, proc_id: ProcId, func: F)
   -> Result<R, Error> where
    F: for<'b> FnOnce(&'b mut Owned<T>) -> R, R: Copy
```

```
fn each<F>(&self, func: F) where
F: for<'b> Fn(&'b mut Owned<T>)
```

Recall: TOSThreads requires 700 bytes statically allocated in the kernel for each additional thread. At most 22% can be used at any given time.

► Grants require no additional per-thread memory in the kernel

- Only useful memory is dynamically allocated in grants
- Zero wasted memory since it can re-use memory for non-concurrent operations.

## Conclusion

- Resource constraints continue to be a challenge for embedded system designers.
  - Low-power, small form-factors and lower cost
- These limitations should not preclude software abstractions and protections common in general-purpose computers.
- Tock provides both dynamic operation and dependability in resource-constrained settings.
  - Best of all: flexible multiprogramming, isolation, system dependability
- Grants split the kernel heap across processes, allowing dynamic demands for kernel resources despite limited system memory

Buy a Hail! https://tockos.org/hardware/hail